

1. The 0th order Bessel function can be computed by adding terms of the power series given as formula 9.1.10 in Abramowitz & Stegun (p. 360).

- a. Write a program to compute the series efficiently by directly adding the series, in the range $0 \leq x \leq 3$. The result should have absolute precision of 10^{-8} .
- b. Use Chebyshev approximation to find a polynomial fit on the same range that provides the same absolute precision. (Hint: use the table 22.3 found on p. 795 of A&S.)

Solution:

A program to evaluate Bessel functions by summing the series is shown on the next page.

In the range $0 \leq x \leq 3$ the 9th term of the series

$$J_0(x) = \sum_{n=0}^{\infty} \frac{(-x^2/4)^n}{(n!)^2}$$

is bounded by 1.12×10^{-8} . If we wish, therefore, we can approximate

$$\xi^9 \stackrel{df}{=} \left(x^2/9\right)^9$$

by

$$\frac{1}{256} [126T_1(\xi) + 84T_3(\xi) + 36T_5(\xi) + 9T_7(\xi)]$$

and ξ^8 by

$$\frac{1}{128} [35T_0(\xi) + 56T_2(\xi) + 28T_4(\xi) + 8T_6(\xi)].$$

The errors from dropping T_8 and T_9 are respectively bounded by

$$\frac{(9/4)^8}{128 \times (8!)^2} \approx 3 \times 10^{-9}$$

and

$$\frac{(9/4)^9}{256 \times (9!)^2} \approx 4 \times 10^{-11}.$$

We cannot economize further since Chebyshev-approximating the term in ξ^7 gives an error of about 2×10^{-7} . Rearranging terms we find

$$\begin{aligned} J_0 \approx & 0.99999999684 - 2.25\xi \\ & + 1.26562510101\xi^2 \\ & - 0.3164062500\xi^3 \\ & + 0.04449412386\xi^4 \\ & - 0.0040045166\xi^5 \\ & + 0.00025109036\xi^6 \\ & - 0.00001149255\xi^7 \end{aligned}$$

Some tests are:

0.0e0	j0	f.	1.0000000	ok
1.0e0	j0	f.	.76519768	ok
2.0e0	j0	f.	.22389078	ok
2.4e0	j0	f.	.00250768	ok
2.5e0	j0	f.	-.04838378	ok
3.0e0	j0	f.	-.26005195	ok

```

\ 0'th order Bessel function by series evaluation
\ precision set to 1.E-8

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\ This is an ANS Forth program requiring the
\ FLOAT wordset.
\
\ Environmental dependences:
\ Assumes independent floating point stack

MARKER -bes0

include ftran201.f

FVARIABLE xx
FVARIABLE sum
FVARIABLE term

: LessThan ( f: a b --) ( -- true if a < b)
  F< ;

: done? f" LessThan(abs(term), 1e-10)" ;

: J0 ( f: x -- J0)
  1 LOCALS| n | \ initialize index
  f" xx=" \ store variable in xx

  f" xx = - (xx/2)^2" \ initialize
  f" sum = 1"
  f" term = xx"

  BEGIN \ loop until |term| < 1e-10
    n 1+ TO n \ increment index
    f" sum = sum + term" \ update sum
    f" term * xx" \ update term
    n S>F f^2 F/
    f" term=" \ and save it

    ( cr n . f" term" f. f" sum" f. ) \ debugging code

  done? UNTIL

  f" sum + term"
;

```

2. The subroutine at the right is intended to solve linear equations with tri-diagonal matrices. It has neither comments nor documentation. Without looking at the “Numerical Recipes” book, analyze the subroutine and provide

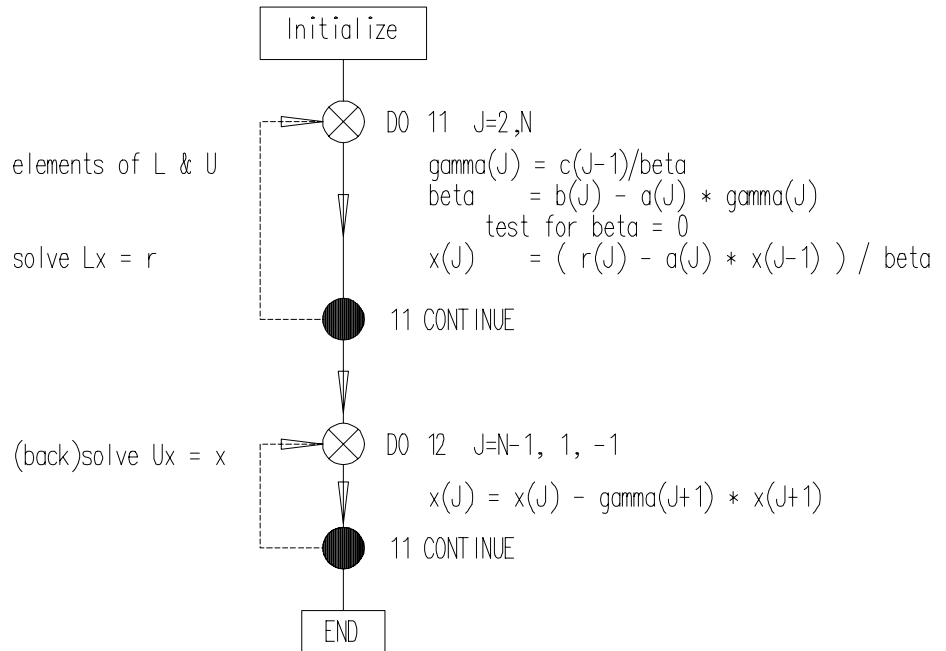
- a) a flow chart;
- b) appropriate comments and documentation in the program;
- c) a working and tested version in your own favorite programming language (different from the above).

```

SUBROUTINE TRIDAG(A,B,C,R,U,N)
PARAMETER (NMAX=100)
DIMENSION GAM(NMAX),A(N),B(N),C(N),R(N),U(N)
IF(B(1).EQ.0.)PAUSE ' Initialize
BET=B(1)
U(1)=R(1)/BET
DO 11 J=2,N ' get elts of L & U
  GAM(J)=C(J-1)/BET
  BET=B(J)-A(J)*GAM(J)
  IF(BET.EQ.0.)PAUSE
  U(J)=(R(J)-A(J)*U(J-1))/BET ' solve Lx=r
11 CONTINUE
DO 12 J=N-1,1,-1 ' solve Ux=x
  U(J)=U(J)-GAM(J+1)*U(J+1)
12 CONTINUE
RETURN
END
    
```

Solution:

Comments and flow chart as shown; program on the Web page.



3. Write a brief program to solve linear systems of equations using the pivotal elimination method (that is, the sequence TRIANGULARIZE \rightarrow BACKSOLVE). Use it to solve the test case

$$\begin{pmatrix} 3 & 2 & 4 & 6 \\ -5 & 5 & 7 & -6 \\ 0 & 18 & 0 & -8 \\ -9 & 9 & 1 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 1 \\ 3 \\ 5 \\ -7 \end{pmatrix}$$

to 6 significant figures of precision. If you wish you may use Maple® or Mathematica® or Matlab® to check your work—or else multiply the matrix by the solution and see whether you obtain the inhomogeneous term.

Solution:

For a program in Forth, see the section “Forth linear equations, etc.” Runs fine under Win32Forth and GForth. (ANS Forth systems.)

The numerical results are

```
a{{ v{ }}solve
x(0 )= .795035
x(1 )= .025639
x(2 )= .491872
x(3 )= -.567312 ok
a{{ Iperm{ .M
1.00000 -1.00000 -.111111 -.111111
.000000 1.00000 .000000 -.444444
.000000 .000000 1.00000 -1.01724
.000000 .000000 .000000 1.00000 ok
```

4. Write a short, efficient program to perform bit-reversal of numbers from 0 to 2^{k-1} where $k=5, 6, \dots, 10$. Make sure you also create a tool to test your answer by displaying the input and output in binary format, in a field of k digits, including the leading 0's. (That is, if the standard integer is 32 bits, there should be $32 - k$ leading 0's.)

Solution:

Here are 3 example programs:

Fortran:

```
INTEGER FUNCTION STIB(N,K)
C
C Return the bit-reversed integer corresponding to the
C K least significant (rightmost) bits of N.
C Algorithm:
C Shift STIB one bit to the left and add the LSB of N.
C Shift N one bit to the right. Repeat K times.
C
C initialize result
C STIB=0
C repeat K times
C initialize result
C STIB=0
C repeat K times
C DO 10 I=1,K
C NP = N/2 right-shift 1 bit
C STIB=N+2*(STIB-NP) left-shift old result, add LSB of N
C N=NP replace N by right-shifted version
10 CONTINUE
RETURN
END
```

C:

```
int stib( int n, int k ) // reverse bit order of #
{
    int i;
    int j=0; // initialize answer
    for( i=1; i<=k; i++ ) // loop k times
    {
        j=2*j+(n && 1); // left-shift result and add 1's bit of n
        n=n/2; // shift n one place to right
    }
    return j;
}
```

Forth:

```
: stib ( n w -- n' ) \ reverse bit order of n in [0,2^w-1]
LOCALS| w n | \ def local variable names
0 ( -- 0) \ n' = 0 is top of stack
w 0 DO ( n' ) \ loop w times
2* \ left-shift n' 1 place
n 1 AND \ get 1's bit of n
+ \ add to n'
n 2/ TO n \ right-shift n -- equiv to n = n/2
LOOP \ end loop
;
```

(see next page for display routine)

Display subroutine in Forth:

```
\ display an integer in binary form, with leading 0's
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\ -----

\ This is an ANS Forth program requiring the CORE wordset

MARKER -bdisp

BL PARSE (D.) DUP PAD C! PAD CHAR+ SWAP CHARS MOVE PAD FIND
NIP 0=
[IF] : (D.) TUCK DABS <# #S ROT SIGN #> ; [THEN]

: ZEROS ( w --)
  0 ?DO [CHAR] 0 EMIT LOOP ;

: 0.R ( n w --) \ display a number right-justified in
                \ a field of width w, with leading 0's
  >R 0 (D.) R>
  OVER - ZEROS TYPE ;

: 0.B ( n w --)
  BINARY
  CR 0.R
  DECIMAL ;

\ Example
\
\ FLOAD C:\WIN32FOR\JVN_PROGS\E_DISP.F ok
\ FLOAD C:\WIN32FOR\JVN_PROGS\STIB1.F ok
\ 17 14 2dup cr 0.b tuck stib swap cr 0.b
\
\ 00000000010001
\
\ 10001000000000 ok
```